Pranking with Ranking

Koby Crammer and Yoram Singer School of Computer Science & Engineering The Hebrew University, Jerusalem 91904, Israel fkobi
s,singerg
s.huji.a
.il

Abstra
t

We discuss the problem of ranking instances. In our framework each instance is associated with a rank or a rating, which is an integer from 1 to k . Our goal is to find a rank-prediction rule that assigns ea
h instan
e a rank whi
h is as lose as possible to the instance's true rank. We describe a simple and efficient online algorithm, analyze its performan
e in the mistake bound model, and prove its orre
tness. We des
ribe two sets of experiments, with synthetic data and with the EachMovie dataset for collaborative filtering. In the experiments we performed, our algorithm outperforms online algorithms for regression and classification applied to ranking.

1 Introdu
tion

The ranking problem we discuss in this paper shares common properties with both classification and regression problems. As in classification problems the goal is to assign one of ^k possible labels to a new instan
e. Similar to regression problems, the set of k labels is structured as there is a total order relation between the labels. We refer to the labels as ranks and without loss of generality assume that the ranks constitute the set $\{1, 2, \ldots, k\}$. Settings in which it is natural to rank or rate instan
es rather than lassify are ommon in tasks su
h as information retrieval and collaborative filtering. We use the latter as our running example. In collaborative filtering the goal is to predict a user's rating on new items such as books or movies given the user's past ratings of the similar items. The goal is to determine whether a movie fan will like a new movie and to what degree, whi
h is expressed as a rank. An example for possible ratings might be, run-to-see, very-good, good, only-if-you-must, and do-not-bother. While the different ratings carry meaningful semantics, from a learning-theoretic point of view we model the ratings as a totally ordered set (whose size is 5 in the example above).

The interest in ordering or ranking of objects is by no means new and is still the source of ongoing research in many fields such mathematical economics, social science, and computer science. Due to lack of space we clearly cannot cover thoroughly previous work related to ranking. For a short overview from a learning-theoreti point of view see $[1]$ and the references therein. One of the main results of $[1]$ unders
ores a omplexity gap between lassi
ation learning and ranking learning. To sidestep the inherent intra
tability problems of ranking learning several approa
hes have been suggested. One possible approach is to cast a ranking problem as a regression problem. Another approa
h is to redu
e a total order into a set of pref-

Figure 1: An Illustration of the update rule.

erences over pairs $[3, 5]$. The first case imposes a metric on the set of ranking rules which might not be realistic, while the second approach is time consuming since it requires increasing the sample size from n to $O(n)$.

In this paper we consider an alternative approach that directly maintains a totally ordered set via projections. Our starting point is similar to that of Herbrich et. al [5] in the sense that we project each instance into the reals. However, our work then deviates and operates directly on rankings by associating each ranking with distinct sub-interval of the reals and adapting the support of each sub-interval while learning. In the next section we describe a simple and efficient online algorithm that manipulates concurrently the direction onto which we project the instances and the division into sub-intervals. In Se
. 3 we prove the orre
tness of the algorithm and analyze its performance in the mistake bound model. We describe in Sec. 4 experiments that compare the algorithm to online algorithms for classification and regression applied to ranking whi
h demonstrate the merits of our approa
h.

$\overline{2}$ 2 The PRank Algorithm

This paper focuses on online algorithms for ranking instances. We are given a sequence $(\mathbf{x}^*, y^*), \ldots, (\mathbf{x}^*, y^*), \ldots$ of instance-rank pairs. Each instance \mathbf{x}^* is in \mathbb{R}^+ and its corresponding rank y^\perp is an element from finite set y with a total order relation. We assume without loss of generality that $\mathcal{Y} = \{1, 2, ..., k\}$ with ">" as the order relation. The total order over the set Y induces a partial order over the instances in the following natural sense. We say that ${\bf x}^*$ is preferred over ${\bf x}^*$ if $y \geq y$, we also say that **x** and **x** are not comparable if neither $y \geq y$ nor $y \sim y$, we denote this case simply as $y \equiv y$. Note that the induced partial order is of a unique form in which the instances form k equivalence classes which are totally ordered . A *ranking rule H* is a mapping from instances to ranks, H : $\mathbb{R}^n \to V$. The family of ranking rules we discuss in this paper employs a vector $w \in \mathbb{R}^n$ and a set of the set of the state by the problems by the set of the se $\mathbf{b} = \mathbf{b}$, $\mathbf{b} = \mathbf{b}$ the vertex of the vector contracts except \mathbf{c} , \mathbf{b} and the vector \mathbf{c} new instance x the ranking rule first computes the inner-product between w and x. The predicted rank is then defined to be the index of the first (smallest) threshold b_r for which $\mathbf{w} \cdot \mathbf{x} < b_r$. This type of ranking rules divide the space into parallel equally-ranked regions: all the instances that satisfy $b_{r-1} < w \cdot x < b_r$ are assigned the same rank r . Formally, given a ranking rule defined by w and b the predicted rank of an instan
e ^x is, H(x) ⁼ minr2f1;:::;kgfr : ^w ^x br < 0g. Note that the above minimum is always well defined since since \mathbb{P}_k = 1.

The analysis that we use in this paper is based on the mistake bound model for online learning. The algorithm we describe works in rounds. On round t the learning algorithm gets an instance ${\bf x}$. Given ${\bf x}$, the algorithm outputs a rank, $y^-=\min_r \{r:$ $\mathbf{w} \cdot \mathbf{x} - \theta_r \leq \theta$. It then receives the correct rank y and updates its ranking rule by modifying **w** and **p**. We say that our algorithm made a ranking mistake if $y^* \neq y^*$.

For a discussion of this type of partial orders see $|b|$.

Imultanze: Set $w^* = 0$, $v_1^*, \ldots, v_{k-1}^* = 0, v_k^* = \infty$. Loop: For ^t = 1; 2; : : : ; T \bullet Get a new rank-value $\mathbf{x}^* \in \mathbb{R}^n$. • Predict $y_{\ell} = \min_{r \in \{1,...,k\}} \{r : \mathbf{w} \mid \mathbf{x} = v_r \leq 0\}.$ $\bullet\,$ Get a new label y . • $\mathbf{u} \quad y^+ \neq y^+$ update wthousand we set wthere $\mathbf{w}^+ \cdot \mathbf{v}^+ = \mathbf{w}^+$, $\forall r : b_r^+ \cdot \mathbf{v}^+ = b_r^+$ r) : 1. For $r = 1, ..., k-1$: If $y \leq r$ Then $y_r = -1$
Else $y_r^t = 1$. 2. For $r = 1, ..., k - 1$: If $(\mathbf{W} \cdot \mathbf{X} - \theta_r) y_r \leq 0$ in entrarching Lise $\tau_r = 0$. 3. Update $\mathbf{w}^{t+1} \leftarrow \mathbf{w}^t + (\sum_r \tau_r^t) \mathbf{x}^t$. For $r = 1, \ldots, \kappa - 1$ update: $v_r^+ \leftarrow v_r^- - \tau_r^-$ **Output :** $H(\mathbf{X}) = \min_{r \in \{1,...,k\}} \{r : \mathbf{W}^T \times \mathbf{X} - \theta_r^T \times \mathbf{U}\}.$

Figure 2: The PRank algorithm.

We wish to make the predicted rank as close as possible to the true rank. Formally, the goal of the learning algorithm is to minimize the ranking-loss which is defined to be the number of thresholds between the true rank and the predicted rank. Using the representation of ranks as integers in $\{1 \dots k\}$, the ranking-loss after T rounds is equal to the accumulated difference between the predicted and true rank-values, $\sum_{t=1}^{T} |\hat{y}^t - y^t|$. The algorithm we describe updates its ranking rule only on rounds on whi
h it made ranking mistakes. Su
h algorithms are alled onservative.

We now describe the update rule of the algorithm which is motivated by the pereptron algorithm for lassi
ation and hen
e we all it the PRank algorithm (for Perceptron Ranking). For simplicity, we omit the index of the round when referring to an input instance-rank pair (x, y) and the ranking rule w and b. Since $\mathbf{1} = \mathbf{2} = \mathbf{3}$ if $\mathbf{4} = \mathbf{1}$ is the predicted rank is r and we have the set of the set of the set of the set of the above interest the above interest the above interest ties by expanding the rank y into into $k-1$ virtual variables y_1, \ldots, y_{k-1} . We set $y_r = +1$ for the case $\mathbf{w} \cdot \mathbf{x} > b_r$ and $y_r = -1$ for $\mathbf{w} \cdot \mathbf{x} < b_r$. Put another way, a rank value y induces the vector $(y_1, \ldots, y_{k-1}) = (+1, \ldots, +1, -1, \ldots, -1)$ where the maximal index r for which $y_r = +1$ is $y-1$. Thus, the prediction of a ranking rule is correct if $y_r(\mathbf{w} \cdot \mathbf{x} - b_r) > 0$ for all r. If the algorithm makes a mistake by ranking **x** as \hat{y} instead of y then there is at least one threshold, indexed r, for which the value of **w** x is on the wrong side of b_r , i.e. $y_r(\mathbf{w} \cdot \mathbf{x} - b_r) \leq 0$. To correct the mistake, we need to "move" the values of $w \cdot x$ and b_r toward each other. We do so by modifying only the values of the b_r 's for which $y_r(\mathbf{w}\cdot \mathbf{x}-b_r) < 0$ and replace them with $b_r - y_r$. We also replace the value of **w** with $\mathbf{w} + (\sum y_r)\mathbf{x}$ where the sum is taken over the indices r for which there was a prediction error, i.e., $y_r(\mathbf{w} \cdot \mathbf{x} - b_r) \leq 0$.

An illustration of the update rule is given in Fig 1. In the example, we used the set $\mathcal{Y} = \{1 \dots 5\}$. (Note that $b_5 = \infty$ is omitted from all the plots in Fig 1.) The correct rank of the instance is $y = 4$, and thus the value of $\mathbf{w} \cdot \mathbf{x}$ should fall in the fourth interval, between b_3 and b_4 . However, in the illustration the value of $\mathbf{w} \cdot \mathbf{x}$ fell below b_1 and the predicted rank is $\hat{y} = 1$. The threshold values b_1 , b_2 and b_3 are e of the since the since the since since \sim is higher the value of the value of the value of the value of th b₁ -1, b₂ -1 and b₃ -1. It also modifies w to be w+3x since $\sum_{r: y_r(w \cdot x - b_r) < 0} y_r = 3$. I mus, the inner-product **w** x increases by $3||x||^2$. This update is illustrated at the middle plot of Fig. 1. The updated prediction rule is sketched on the right hand side of Fig. 1. Note that after the update, the predicted rank of **x** is $\hat{y} = 3$ which is closer to the true rank $y = 4$. The pseudocode of algorithm is given in Fig 2.

To conclude this section we like to note that PRank can be straightforwardly combined with Mercer kernels [8] and voting techniques [4] often used for improving the performan
e of margin lassiers in bat
h and online settings.

3 Analysis

Before we prove the mistake bound of the algorithm we first show that it maintains a onsistent hypothesis in the sense that it preserves the orre
t order of the be derived by the algorithm along its run, $(\mathbf{w}^1, \mathbf{b}^1)$..., $(\mathbf{w}^{T+1}, \mathbf{b}^{T+1})$ we have that $v_r \leq \ldots \leq v_{k-1}$ for all t. Since the initialization of the thresholds is such that $v_1 \le v_2 \le \ldots \le v_{k-1}$, then it sumess to show that the claim holds inductively. For simplicity, we write the updating rule of PRank in an alternative form. Let $\llbracket \pi \rrbracket$ be I if the predicate π holds and 0 otherwise. We now rewrite the value of τ_r^* (from Fig. 2) as $\tau_r = y_r \llbracket (\mathbf{w} \cdot \mathbf{x} - v_r) y_r \leq 0 \rrbracket$. Note that the values of v_r are integers for all r and t since for all r we initialize $v_r^r = 0$, and $v_r^r - v_r \in \{-1, 0, \pm 1\}$.

Lemma 1 (Order Preservation) Let w^2 and p^2 be the current ranking rate, where $v_1 \leq \ldots \leq v_{k-1}$, and let $(\mathbf{x}^\star, y^\star)$ be an instance-rank pair fed to Prank on round t. Denote by wt+1 and bt+1 the resulting ranking rule after the update of PRank, then $b_1 \leq \cdots \leq b_{k-1}$.

Proof: In order to show that PRank maintains the order of the thresholds we use the definition of the algorithm for y_r , namely we define $y_r = +1$ for $r < y^+$ and $y_r^* = -1$ for $r \geq y^*$ we now prove that $b_{r+1} \geq b_r^{r+1}$ for all r by showing that

$$
b_{r+1}^t - b_r^t \ge y_{r+1}^t \left[(\mathbf{w}^t \cdot \mathbf{x}^t - b_{r+1}^t) y_{r+1}^t \le 0 \right] - y_r^t \left[(\mathbf{w}^t \cdot \mathbf{x}^t - b_r^t) y_r^t \le 0 \right], \quad (1)
$$

which we obtain by substituting the values of **b**^{the}. Since $b_{r+1}^* \leq b_r^*$ and b_r^* , $b_{r+1}^* \in \mathbb{Z}$ we get that the value of $b_{r+1}^* - b_r^*$ on the left hand side of Eq. (1) is a non-negative integer. Recall that $y_r = 1$ if $y \geq r$ and $y_r = -1$ otherwise, and therefore,
 $y_{r+1}^t \leq y_r^t$. We now analyze two cases. We first consider the case $y_{r+1}^t \neq y_r^t$ which implies that $y_{r+1} = -1$, $y_r = +1$. In this case, the right hand-side of Eq. (1) is at most zero, and the claim trivially holds. The other case is when $y_{r+1} = y_r$. Here we give that the value of the right hand-side Eq. (1) therefore Eq. (1) there is a set α have to consider only the case where $v_r = v_{r+1}$ and $y_{r+1} = y_r$. But given these two conditions we have that y_{r+1}^* $\|(\mathbf{w} \cdot \mathbf{x} - v_{r+1}^*)y_{r+1}^* \leq v\|$ and y_r^* $\|(\mathbf{w} \cdot \mathbf{x} - v_r^*)y_r^* \leq v\|$ are equal. The right mass side of Eq. (1) is now zero and the inequality method with equality.

In order to simplify the analysis of the algorithm we introdu
e the following notation. Given a hyperplane w and a set of $k-1$ thresholds **b** we denote by $\mathbf{v} \in \mathbb{R}^{n+k-1}$ the vector which is a concatenation of **w** and **b** that is $\mathbf{v} = (\mathbf{w}, \mathbf{b})$. For brevity we refer to the vector **v** as a *ranking rule*. Given two vectors $\mathbf{v} = (\mathbf{w}, \mathbf{b})$ and $\mathbf{v} = (\mathbf{w}, \mathbf{b})$ we have $\mathbf{v} - \mathbf{v} = \mathbf{w} - \mathbf{w} + \mathbf{b} - \mathbf{b}$ and $\|\mathbf{v}\| = \|\mathbf{w}\| = + \|\mathbf{b}\|$.

Theorem 2 (MISTAKE DOUND) Let $(\mathbf{x}^*, y^*, \dots, (\mathbf{x}^*, y^*)$ be an input sequence for

PRank where $\mathbf{x}^t \in \mathbb{R}^n$ and $y^t \in \{1 \dots k\}$. Denote by $R^2 = \max_t ||\mathbf{x}^t||^2$. Assume

that there is a ranking rule $\mathbf{v}^* = (\mathbf{w$ classifies the entire sequence correctly with margin $\gamma = \min_{r,t} \{ (\mathbf{w} \cdot \mathbf{x} - \theta_r) y_r \} \geq 0.$ Then, the rank loss of the algorithm $\sum_{t=1}^{I} |\hat{y}^t - y^t|$, is at most $(k-1)(R^2 + 1)/\gamma^2$.

Proof: Let us fix an example (X, y') which the algorithm received on round t. By definition the algorithm ranked the example using the ranking rule \mathbf{v}^\top which is composed of w^* and the thresholds p^* . Similarly, we denote by v^* - the updated rule $(\mathbf{w}^{t+1}, \mathbf{b}^{t+1})$ after round t. That is, $\mathbf{w}^{t+1} = \mathbf{w}^t + (\sum_r \tau_r^t) \mathbf{x}^t$ and $b_r^{t+1} = b_r^t - \tau_r^t$
for $r = 1, 2, ..., k-1$. Let us denote by $n^t = |\hat{y}^t - y^t|$ the difference between the true
rank and the p rank and the predicted rank. It is straightforward to verify that $n^t = \sum_r |\tau_r^t|$. Note that if there wasn't a ranking mistake on round t then $\tau_r^t = 0$ for $r = 1, \ldots, k-1$, and thus also $n^t = 0$. To prove the theorem we bound $\sum_t n^t$ from above by bounding $\mathbf{v}^* \cdot \mathbf{v}^{t+1}$. Substituting the values of \mathbf{w}^{t+1} and \mathbf{b}^{t+1} we get,

$$
\mathbf{v}^* \cdot \mathbf{v}^{t+1} = \mathbf{v}^* \cdot \mathbf{v}^t + \sum_{r=1}^{k-1} \tau_r^t \left(\mathbf{w}^* \cdot \mathbf{x}^t - b_r^* \right)
$$
 (2)

We further bound the right term by considering two cases. Using the definition of τ_r from the pseudocode in Fig. 2 we need to analyze two cases. If $(\mathbf{w} \cdot \mathbf{x} - \theta_r)y_r$ $r =$ then $\tau_r^{\star} = y_r^{\star}$. Using the assumption that ${\bf v}$ -ranks the data correctly with a margin of at least γ we get that $\tau_r^{\text{f}}(\mathbf{w} \cdot \mathbf{x}^t - v_r) \geq \gamma$. For the other case for which
 $(\mathbf{w}^t \cdot \mathbf{x}^t - b_r^t)y_r^t > 0$ we have $\tau_r^t = 0$ and thus $\tau_r^t(\mathbf{w}^* \cdot \mathbf{x}^t - b_r^t) = 0$. Summing now over r we get,

$$
\sum_{r=1}^{k-1} \tau_r^t \left(\mathbf{w}^* \cdot \mathbf{x}^t - b_r^* \right) \ge n^t \gamma \tag{3}
$$

Combining Eq. (2) and Eq. (3) we get $\mathbf{v} \cdot \mathbf{v} = 0$ $\mathbf{v} \cdot \mathbf{v} + n\gamma$. Unioning the sum, we get that after T rounds the algorithm satisfies, $\mathbf{v}^* \cdot \mathbf{v}^{T+1} \ge \sum_t n^t \gamma = \gamma \sum_t n^t$. Plugging this result into Ca $(v^{T+1} \cdot v^*)^2$ and using the assumption that v^* is of a unit norm we get the lower bound, $\|\mathbf{v}^{T+1}\|^2 \geq (\sum_t n^t)^2 \gamma^2$.

Next, we bound the norm of v from above. As before, assume that an example (x^*, y^*) was ranked using the ranking rule v^* and denote by v^* + the ranking rule after the round. We now expand the values of $\mathbf{w}=\tan\mathbf{u}$ better in the norm of $\mathbf{v}=\tan\alpha$ $\text{get},\, \|\mathbf{v}^{t+1}\|^2 = \|\mathbf{w}^t\|^2 + \|\mathbf{b}^t\|^2 + 2\sum_r \tau_r^t\,(\mathbf{w}^t\cdot\mathbf{x}^t - b_r^t) + (\sum_r \tau_r^t)^2 \|\mathbf{x}^t\|^2 + \sum_r \left(\tau_r^t\right)^2.$ Since $\tau_r^t \in \{-1,0,+1\}$ we have that $(\sum_r \tau_r^t)^2 \leq (n^t)^2$ and $\sum_r (\tau_r^t)^2 = n^t$ and we therefore get,

$$
\|\mathbf{v}^{t+1}\|^2 \le \|\mathbf{v}^t\|^2 + 2\sum_r \tau_r^t \left(\mathbf{w}^t \cdot \mathbf{x}^t - b_r^t\right) + (n^t)^2 \|\mathbf{x}^t\|^2 + n^t. \tag{4}
$$

We further develop the se
ond term using the update rule of the algorithm and get,

$$
\sum_{r} \tau_r^t \left(\mathbf{w}^t \cdot \mathbf{x}^t - b_r^t \right) = \sum_{r} \left[(\mathbf{w}^t \cdot \mathbf{x}^t - b_r^t) y_r^t \le 0 \right] \left((\mathbf{w}^t \cdot \mathbf{x}^t - b_r^t) y_r^t \right) \le 0 \ . \tag{5}
$$

Plugging Eq. (5) into Eq. (4) and using the bound $||\mathbf{x}^*||^2 \leq \mathbf{h}^{-1}$ we get that $||\mathbf{v}^{t+1}||^2 \leq ||\mathbf{v}^t||^2 + (n^t)^2 R^2 + n^t$. Thus, the ranking rule we obtain after T rounds of the algorithm satisfies the upper finally get, -^P -

$$
\sum_{t} n^{t} \le \frac{R^{2} \left[\sum_{t} (n^{t})^{2}\right] / \left[\sum_{t} n^{t}\right] + 1}{\gamma^{2}} \tag{6}
$$

By definition, n^t is at most $k-1$, which implies that $\sum_{t} (n^t)^2 \leq \sum_{t} n^t (k-1) =$ $(k-1)\sum_t n^t$. Using this inequality in Eq. (6) we get the desired bound, $\sum_{t=1}^T |\hat{y}^t |y^t| = \sum_{t=1}^T n^t \leq [(k-1)R^2+1]/\gamma^2 \leq [(k-1)(R^2+1)]/\gamma^2$.

Figure 3: Comparison of the time-averaged ranking-loss of PRank, WH, and MCP on syntheti data (left). Comparison of the time-averaged ranking-loss of PRank, WH, and MCP on the EachMovie dataset using viewers who rated and at least 200 movies (middle) and at least 100 movies (right).

4 Experiments

In this section we describe experiments we performed that compared PRank with two other online learning algorithms applied to ranking: a multi
lass generalization of the perceptron algorithm $[2]$, denoted MCP, and the Widrow-Hoff $[9]$ algorithm for online regression learning which we denote by WH. For WH we fixed its learning rate to a onstant value. The hypotheses the three algorithms maintain share similarities but are different in their complexity: PRank maintains a vector \bf{w} of dimension n and a vector of $k-1$ modifiable thresholds **b**, totaling $n+k-1$ parameters; MCP maintains k prototypes which are vectors of size n , yielding kn parameters; WH maintains a single vector \bf{w} of size n. Therefore, MCP builds the most omplex hypothesis of the three while WH builds the simplest.

Due to the lack of space, we only describe two sets of experiments with two different datasets. The dataset used in the first experiment is synthetic and was generated in a similar way to the dataset used by Herbrich et. al. $[5]$. We first generated random points $\mathbf{x} = (x_1, x_2)$ uniformly at random from the unit square $[0, 1]$. Each point was assigned a rank y from the set $\{1, \ldots, 5\}$ according to the following ranking rule, $y = \max_r \{r : 10((x_1 - 0.5)(x_2 - 0.5)) + \xi > b_r\}$ where $\mathbf{b} = (-\infty, -1, -0.1, 0.25, 1)$ and ξ is a normally distributed noise of a zero mean and a standard deviation of 0.125. We generated 100 sequen
es of instan
e-rank pairs ea
h of length 7000. We fed the sequences to the three algorithms and obtained a prediction for each instan
e. We onverted the real-valued predi
tions of WH into ranks by rounding each prediction to its closest rank value. As in [5] we used a non-homogeneous polynomial of degree $\mathbf{z}, \, \mathbf{A} \, (\mathbf{x}_1, \mathbf{x}_2) = ((\mathbf{x}_1 \cdot \mathbf{x}_2) + 1)^2$ as the inner-product operation between ea
h input instan
e and the hyperplanes the three algorithms maintain. At each time step, we computed for each algorithm the accumulated ranking-loss normalized by the instantaneous sequen
e length. Formally, the time-averaged loss after T rounds is, $(1/T) \sum_{t}^{T} |\hat{y}^{t} - y^{t}|$. We computed these losses for $T = 1, ..., 7000$. To in
rease the statisti
al signi
an
e of the results we repeated the pro
ess 100 times, picking a new random instance-rank sequence of length 7,000 each time, and averaging the instantaneous losses across the 100 runs. The results are depicted on the left hand side of Fig. 3. The 95% confidence intervals are smaller then the symbols used in the plot. In this experiment the performance of MPC is constantly worse than the performance of WH and PRank. WH initially suffers the smallest instantaneous loss but after about 500 rounds PRank a
hieves the best performan
e and eventually the number of ranking mistakes that PRank suffers is significantly lower than both WH and MPC.

In the second set of experiments we used the EachMovie dataset [7]. This dataset is used for ollaborative ltering tasks and ontains ratings of movies provided by 61; 265 people. Ea
h person in the dataset viewed a subset of movies from a collection of 1623 titles. Each viewer rated each movie that she saw using one of 6 possible ratings: $0, 0.2, 0.4, 0.6, 0.8, 1$. We chose subsets of people who viewed a significant amount of movies extracting for evaluation people who have rated at least 100 movies. There were 7; 542 su
h viewers. We hose at random one person among these viewers and set the person's ratings to be the target rank. We used the ratings of all the rest of the people who viewed enough movies as features. Thus, the goal is to learn to predict the "taste" of a random user using the user's past ratings as a feedba
k and the ratings of fellow viewers as features. The predi
tion rule asso
iates a weight with ea
h fellow viewer an therefore an be seen as learning correlations between the tastes of different viewers. Next, we subtracted 0.5 from each rating and therefore the possible ratings are $-0.5, -0.3, -0.1, 0.1, 0.3, 0.5$. This linear transformation enabled us to assign a value of zero to movies whi
h have not been rated. We fed these feature-rank pairs one at a time, in an online fashion. Sin
e we pi
ked viewer who rated at least 100 movies, we were able to perform at least 100 rounds of online predictions and updates. We repeated this experiment 500 times, hoosing ea
h time a random viewer for the target rank. The results are shown on the right hand-side of Fig. 3. The error bars in the plot indicate 95% condfidence levels. We repeated the experiment using viewers who have seen at least 200 movies. (There were 1802 su
h viewers.) The results of this experiment are shown in the middle plot of Fig. 3. Along the entire run of the algorithms, PRank is significantly better than WH, and consistently better than the multiclass per
eptron algorithm, although the latter employs a bigger hypothesis.

Finally, we have also evaluated the performan
e of PRank in a bat
h setting, using the experimental setup of $[5]$. In this experiment, we ran PRank over the training data as an online algorithm and used its last hypothesis to rank unseen test data. Here as well PRank came out first, outperforming all the algorithms described in [5].

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